The Cycle Structure of Quandles

Naqeeb ur Rehman

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References:

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- S. Guest and P. Spiga, Finite primitive groups and regular orbits of group elements, Trans. Amer. Math. Soc., (2016).

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A rack X is called **indecomposable or connected** if Inn(X) acts transitively on X.

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• **Conjugation Quandle** on a group *G* with $x \triangleright y = xyx^{-1}$.

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- Let A be an abelian group, α ∈ Aut(A) and 1 = id_A. Then we have a quandle structure on A, defined by:

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Let G be a group, α ∈ Aut(G), and H be a subgroup of the fixed points of α in G. Then for any g, f ∈ G, the quandle structure on G/H is defined by:

$$gH \triangleright fH = g\alpha(g^{-1}f)H.$$

This quandle is known as **coset quandle** (G, H, α) .

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Remarks

• Since
$$(x \triangleright y) \triangleright (x \triangleright z) = x \triangleright (y \triangleright z)$$
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 Hayashi called the pattern of any φ_x as the profile of a finite connected rack X for short.

Notation

We write the profile of a finite connected rack X as: $\begin{aligned} & Profile(X) = 1^{m_0} \ell_1^{m_1} \ell_2^{m_2} \dots \ell_k^{m_k}, \\ \text{where } 1 < \ell_1 < \ell_2 < \dots < \ell_k, \text{ and } m_0, m_1, \dots, m_k \text{ are the multiplicities of } 1, \ell_1, \ell_2, \dots, \ell_k, \text{ respectively.} \end{aligned}$

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Let X be a finite connected quandle with

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Then $\ell_i | \ell_k$ (i.e., ℓ_i divides ℓ_k) for any integer *i* with $1 \le i \le k - 1$.

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Example

$$Profile(SmallQuandle(42,7)) = 1^2.2^2.3^4.6^4.$$
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Hence, $Profile(\mathbb{D}_n) = 1^{m_0} 2^{m_1}$.

L. Vendramin calculated all connected quandles of size n ≤ 47. These small quandles are included in a GAP package called **Rig** (Racks in gap) as: **SmallQuandle(n, q(n))**, where q(n) := quandle number of size n.

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- Such a cycle σ is called a **regular cycle or orbit**. For example, $\alpha = (1 \ 2)(3 \ 4 \ 5)(6 \ 7 \ 8 \ 9 \ 10 \ 11)$ has a regular cycle, while $\beta = (1 \ 2)(3 \ 4 \ 5)(6 \ 7 \ 8 \ 9)$ has no regular cycle.
- Recall that ord(α) is the least common multiple (lcm) of the cycle lengths of α. Therefore, if α has a regular cycle σ then all cycle lengths of α divide the l(σ) = ord(α).
- For an affine quandle $Aff(A, \alpha)$ we have:

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If we take x = e, the identity of A, then $\varphi_e(y) = \alpha(y)$. Hence, the cycle structure of φ_e is same as the cycle structure of $\alpha \in Aut(A)$. Now, since an abelian group is nilpotent, the automorphism α has a regular cycle.

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• The results of A. Bors, S. Guest and P. Spiga can be used for case-by-case analysis of some other known families of connected quandles.

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Lemma 1

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Proof.

Since Y is a subrack of X, we conclude that $Y \triangleright Y^c = Y^c$. Let

$$Z = \{y_1 \triangleright (y_2 \triangleright ... \triangleright (y_{n-1} \triangleright y_n)) \mid n \ge 1, y_1, ..., y_n \in Y^c\}.$$

Then $y \triangleright z \in Z$ for all $y \in Y^c$, $z \in Z$ by definition, and $y \triangleright z \in Z$ for all $y \in Y$ by the self-distributivity of \triangleright and the Y-invariance of Y^c . Hence Z is a non-empty X-invariant subset of X, and therefore equal to X since X is connected.

Lemma 2

Let X be a connected rack such that $X = Y \cup Z$, for two subracks Y and Z of X. Then X = Y or X = Z.

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Assume that $X \neq Y$, then X is generated by $Y^c \subseteq Z$, by Lemma 1. Since Z is a subrack of X, one concludes that X = Z.

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Subracks of a Connected Rack

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Assume that $X \neq Y$, then X is generated by $Y^c \subseteq Z$, by Lemma 1. Since Z is a subrack of X, one concludes that X = Z.

Corollary 1

Let X be a connected rack and $x \in X$. Let $p, q \in \mathbb{N}_{\geq 2}$, and $Y = \{y \in X \mid \varphi_x^p(y) = y\}, Z = \{z \in X \mid \varphi_x^q(z) = z\}$. Assume that $X = Y \cup Z$. Then X = Y or X = Z.

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Proposition 1.

There is no finite connected rack X (respectively, quandle) of profile $1^{m_0} \ell_1^{m_1} \ell_2^{m_2} \dots \ell_k^{m_k}$ such that $lcm(\ell_1, \ell_2, \dots, \ell_i)$ and $lcm(\ell_{i+1}, \ell_{i+2}, \dots, \ell_k)$ do not divide each other.

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Proof.

Suppose that there exists a finite connected rack X with given profile. Let $p = lcm(\ell_1, \ell_2, ..., \ell_i)$, and $q = lcm(\ell_{i+1}, \ell_{i+2}, ..., \ell_k)$. Then,

$$Y = \{y \in X \mid \varphi_x^p(y) = y\}, Z = \{z \in X \mid \varphi_x^q(z) = z\}.$$

By the self-distributivity of \triangleright , the sets Y and Z are subracks of X. Then $X = Y \cup Z$ by definition of p and q and, $X \neq Y$ and $X \neq Z$, a contradiction to Corollary 1.

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• By above results, there is no finite connected rack X of profile $1^{m_0} \ell_1^{m_1} \ell_2^{m_2}$ with $\ell_1 \nmid \ell_2$.

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• For the profile $1^{m_0}\ell_1^{m_1}\ell_2^{m_2}\ell_3^{m_3}$, we have the following possibilities.

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- For the profile $1^{m_0}\ell_1^{m_1}\ell_2^{m_2}\ell_3^{m_3}$, we have the following possibilities.
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 $1^{m_0}6.10.15 = 1^{m_0}(2.3)(2.5)(3.5)$ with (p, q, r) = (2, 3, 5).

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Proposition 2

There is no finite connected crossed set with profile $1^{m_0}\ell_1\ell_2\ell_3$, where $(\ell_1, \ell_2, \ell_3) = (pq, pr, qr)$ for pairwise distinct primes p, q, r.

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Suppose that there exists a finite connected crossed set X with given profile. Let $x \in X$, $t \ge 1$,

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X_t is a non-empty subrack of X. In particular,

$$y \rhd (X \setminus X_t) = X \setminus X_t$$
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For all
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Step 3

Let $y \in X'_{pq}$ and $z \in X \setminus X_{pq}$ with $y \triangleright z \neq z$. Let t be the smallest positive integer with $\varphi_y^t(z) = z$, then t = pr or t = qr.

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Step 4

Let
$$y \in X'_{pq}$$
. Then there exist $z \in X'_{pr}$, $f \in X'_{qr}$ such that $z \triangleright f = y$ or $f \triangleright z = y$.

Assume that X'_{pr} and X'_{qr} are subracks of X. Then X'_{pq} is not a subrack of X.

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Step 6

Let $y \in X'_{pq}$. Then there exist $z \in X'_{pr}$ and $f \in X'_{qr}$ with $y \triangleright z \in X'_{qr}$, $y \triangleright f \in X'_{pr}$.

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Step 7

Let $y \in X'_{pq}$. Then $y \triangleright x \in X'_{pq}$.

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Step 8

Let $y \in X'_{pq}$. Let $\varphi_y = \sigma_1 \sigma_2 \sigma_3$ be the decomposition of φ_y into the product of a pq-, pr-, and qr-cycle. Then $supp(\sigma_1) \subseteq X_{pq}$ and $supp(\sigma_2), supp(\sigma_3) \subseteq X'_{pr} \cup X'_{qr}$.

Let $y \in X'_{pq}$ and $z \in X'_{pr}$, then $z \triangleright x \neq x$ and $\varphi_z^{pr}(x) = x$, by Step 7. Therefore $\varphi_{y \triangleright z}^{pr}(y \triangleright x) = y \triangleright x$. Moreover, $y \triangleright x \in X'_{pq}$, by Step 7. Step 1 implies that $y \triangleright z \in X'_{pr} \cup X'_{qr}$. If $y \triangleright z \in X'_{pr}$, then the entries of pr-cycle of $\varphi_{y \triangleright z}$ belong to X_{pr} , by Step 8, in contradiction to $y \triangleright x \in X'_{pq}$. Thus $y \triangleright z \in X'_{qr}$, which implies that $y \triangleright X'_{pr} \subseteq X'_{qr}$ and $y \triangleright X'_{qr} \subseteq X'_{pr}$ by symmetry. This is impossible since $|X'_{pr}| = pr \neq qr = |X'_{qr}|$.

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Remarks

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Remarks

• By Proposition 2, there is no connected crossed set with profile 1^{*m*}₀6.10.15.

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Remarks

- By Proposition 2, there is no connected crossed set with profile 1^m₀6.10.15.
- By Propositions 1, 2, and Corollary 1, Hayashi's conjecture is true for any connected crossed set X with φ_x ∈ Aut(X) such that supp(φ_x) ≤ 31.

Let $y \in X'_{pq}$ and $z \in X'_{pr}$, then $z \triangleright x \neq x$ and $\varphi_z^{pr}(x) = x$, by Step 7. Therefore $\varphi_{y \triangleright z}^{pr}(y \triangleright x) = y \triangleright x$. Moreover, $y \triangleright x \in X'_{pq}$, by Step 7. Step 1 implies that $y \triangleright z \in X'_{pr} \cup X'_{qr}$. If $y \triangleright z \in X'_{pr}$, then the entries of pr-cycle of $\varphi_{y \triangleright z}$ belong to X_{pr} , by Step 8, in contradiction to $y \triangleright x \in X'_{pq}$. Thus $y \triangleright z \in X'_{qr}$, which implies that $y \triangleright X'_{pr} \subseteq X'_{qr}$ and $y \triangleright X'_{qr} \subseteq X'_{pr}$ by symmetry. This is impossible since $|X'_{pr}| = pr \neq qr = |X'_{qr}|$.

Remarks

- By Proposition 2, there is no connected crossed set with profile 1^{m0}6.10.15.
- By Propositions 1, 2, and Corollary 1, Hayashi's conjecture is true for any connected crossed set X with φ_x ∈ Aut(X) such that supp(φ_x) ≤ 31.
- Proposition 2 is a particular case of the following theorem.

Theorem

Let $p_1, p_2, ..., p_r$ be pairwise distinct primes for positive integer r. Let

$$\ell_1 = \prod_{i=1}^r p_i^{a_i}$$
, $\ell_2 = \prod_{i=1}^r p_i^{b_i}$ and $\ell_3 = \prod_{i=1}^r p_i^{c_i}$,

for non-negative integers a_i , b_i and c_i for all $1 \le i \le r$. Let $1 < \ell_1 < \ell_2 < \ell_3$, $\ell_1 \nmid \ell_3$, $\ell_2 \nmid \ell_3$ and $\ell_k \mid lcm(\ell_{k+1}, \ell_{k+2}) \pmod{3}$ for $k \in \{1, 2, 3\}$. Then there is no finite connected crossed set X with profile $1^{m_0}\ell_1\ell_2\ell_3$.
Obstruction on the Profile of a Connected Rack

Theorem

Let $p_1, p_2, ..., p_r$ be pairwise distinct primes for positive integer r. Let

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Questions, comments, suggestions???